

IX. System Models (III)

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Objective

- Introduce formal descriptive notation
 - 🖲 Logic
 - Algebraic





- Use of First Order Logic formula to describe program properties
- Expressions involve variables, numeric constants, functions, predicates. Logical connectives and, or, not, implies, for all, exists
- 🖻 Examples

x > y and y > z implies x>z
for all i in N x[i] > x[i+1]
for all M in N (exist n in N(n>M))

Free variables and bound variables





Logic Specifications – input/output assertions

A property for a program or a subprogram P is specified as a formula of type
 {Pre(i₁,i₂,...,i_n)}
 P

{Post($O_1, O_2, ..., O_m, i_1, i_2, ..., i_n$)}





input/output assertions - example

{true} Ρ $\{(0=i_1 \text{ or } 0=i_2) \text{ and } 0 \ge i_1 \text{ and } 0 \ge i_2\}$ ▼ {n>0} Ρ $\{0 = \sum_{k=1...n} i_k\}$ \checkmark {(i₁>0 **and** i₂>0)} Ρ {(**exists z**₁, z_2 ($i_1 = 0 \times z_1$ and $i_2 = 0 \times z_2$) and not (**exists** $h(i_1=h \times z_1 \text{ and } i_2=h \times z_2)$ and h>0))





Generalising input/output assertions

- We would like to talk of **programs fragments**!
- The language should be able to refer variables defined within the program
- (n > 0) n is a constant value
 procedure search(table: in integer_array; n: in integer; element: in integer; found: out Boolean); {found=(exist i(1<=i<=n and table(i)=element))}</p>



input/output assertions and classes

The language "predicates" over object variables and data structures defining:

Invariants

- Object states
- Methods execution must preserve invariants
 - Examples: set abstraction or ordered set abstraction

Pre- and Post-conditions on each methods

- INV and precondition} program fragment for m {INV and post-condition}
- Logical statements defined by the developer to increase verification power





Example

```
public class Queue {
  Object[] queue;
  int length, head, tail, elementInQueue;
  public Queue(int length) {
    queue = new Object[length];
    this.length=length;this.head=0;this.tail=0;this.elementInOueue=0;
  }
  public void insert(Object o) {
    gueue[head]=o;head++;elementInOueue++;
    if (head==length) head=0;
  }
  public Object remove() {
    Object o = queue[tail];tail++;elementInQueue--;
    if (tail==length) tail=0;
    return o;
  public boolean isEmpty() {
    return (head==tail);
```



Design by Contract

- Engineering principle known as Design by Contract (DbC)
 - Logical statements defined at the interface level
 - Constitute the agreement among the contractor and the client
 - Support in programming languages: Eiffel, JContractor ...
- Design vs. Requirements using logic specifications
 - Which is the universe of discourse?





- Using FOT you specify that any implementation must guarantee that all of the given rules are true
- Logical expressions can be used to analyze system properties applying logical deduction
 - Derive the formula from the specification of the system
- Operational formalisms do not allow to prove properties
- You can discover behaviour but you cannot provide a proof!
- Proving theorems is not a decidable problem within FOT





- Large systems are decomposed into subsystems with welldefined interfaces between these subsystems.
- Specification of subsystem interfaces allows independent development of the different subsystems.
- Interfaces may be defined as abstract data types or object classes.
- The algebraic approach to formal specification is particularly well-suited to interface specification as it is focused on the defined operations in an object.





Specification elements

- Introduction
 - Defines the sort (the type name) and declares other specifications that are used.
- Description
 - Informally describes the operations on the type.
- Signature
 - Defines the syntax of the operations in the interface and their parameters.
- Axioms
 - Defines the operation semantics by defining axioms which characterise behaviour.





Systematic Algebraic Specification

- Algebraic specifications of a system may be developed
 - in a systematic way
 - Specification structuring;
 - Specification naming;
 - Operation selection;
 - Informal operation specification;
 - Syntax definition;
 - Axiom definition.





- Constructor operations. Operations which create entities of the type being specified.
- Inspection operations. Operations which evaluate entities of the type being specified.
- Rule of thumb: identify constructor operations and define inspector operations for each primitive constructor operation.





Operations on a List ADT

- Constructor operations which evaluate to sort List
 Create, Cons and Tail.
- Inspection operations which take sort list as a parameter and return some other sort
 - Head and Length.
- Tail can be defined using the simpler constructors Create and Cons. No need to define Head and Length with Tail.





Operations on a List ADT

LIST(ELEM)

sort List imports INTEGER

Defines a list where elements are added at the end and removed from the front. The operations are Create, which brings an empty list into existence, Cons, which creates a new list with an added member, Length, which evaluates the list size, Head, which evaluates the front element of the list, and Tail, which creates a list by removing the head from its input list. Undefined represents an undefined value of type Elem.

```
Create --> List
Cons(List, Elem) --> List
Head(List) --> Elem
Length(List) --> Integer
Tail(List) --> List
```

```
\begin{aligned} &\text{Head}(\text{Create}) = \text{Undefined exception (empty list)} \\ &\text{Head}(\text{Cons}(L,v)) = \text{if } L = \text{Create then } v \text{ else } \text{Head}(L) \\ &\text{Length}(\text{Create}) = 0 \\ &\text{Length}(\text{Cons}(L,v)) = \text{Length}(L) + 1 \\ &\text{Tail}(\text{Create}) = \text{Create} \\ &\text{Tail}(\text{Cons}(L,v)) = \text{if } L = \text{Create then } \text{Create else } \text{Cons}(\text{Tail}(L),v) \end{aligned}
```



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Operations on a List ADT

- Operations are often specified recursively.
- Tail (Cons (L, v)) = **if** L = Create **then** Create

else Cons (Tail (L), v).

- Cons ([5, 7], 9) = [5, 7, 9]
- Tail ([5, 7, 9]) = Tail (Cons ([5, 7], 9)) =
- Cons (Tail ([5, 7]), 9) = Cons (Tail (Cons ([5], 7)), 9) =
- Cons (Cons (Tail ([5]), 7), 9) =
- Cons (Cons (Tail (Cons ([], 5)), 7), 9) =
- Cons (Cons ([Create], 7), 9) = Cons ([7], 9) = [7, 9]





Algebraic specifications

- Algebraic Specifications can be easily combined
- Reuse of algebras
 - Import allow to reuse
 - Assume allows to rewrite definition





Keynote

- Discussed two formalism for defining a software specification with a descriptive flavour
 - Use of FOT
 - Algebraic specification
- Operational spec better for simulation descriptive for analysis
- Limitations of analysis have been highlighted



